

# Communications over Sparse Channels: Fundamental limits and practical design

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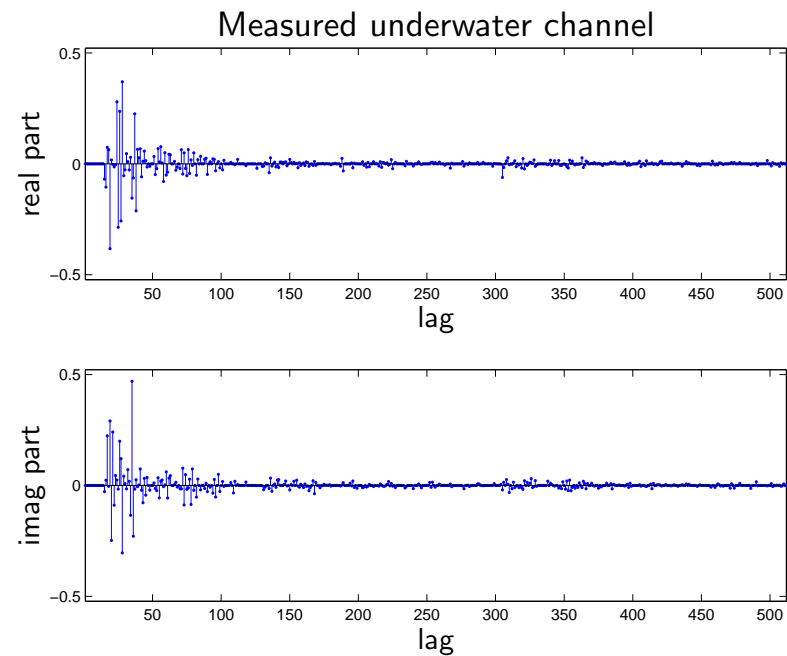
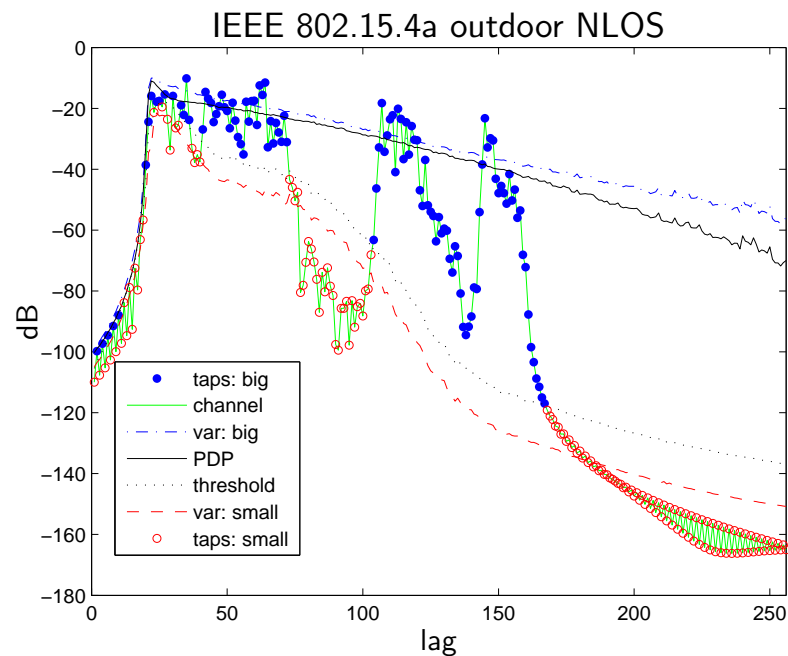
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## Sparse Channels:

- At large communication bandwidths, channel impulse responses are **sparse**.
- Below left shows channel taps  $\mathbf{x} = [x_0, \dots, x_{L-1}]$ , where
  - $x_n = x(nT)$  for bandwidth  $T^{-1} = 256$  MHz,
  - $x(t) = h(t) * p_{RC}(t)$ , and
  - $h(t)$  is generated randomly using 802.15.4a outdoor NLOS specs.



## Simplified Channel Model:

First, let's **simplify** things to talk concretely about sparse channels...

Consider a discrete-time channel that is

- **block-fading** with block size  $N$ ,
- **frequency-selective** with impulse response length  $L$  (where  $L < N$ ),
- **sparse** with  $S$  non-zero complex-Gaussian taps (where  $0 < S \leq L$ ),

where *both the channel coefficients and support are unknown* to the receiver.

Important questions:

1. What is the **capacity** of this channel?
2. How can we build a **practical** comm system that operates near this capacity?

## Noncoherent Capacity of the Sparse Channel:

For the **unknown  $N$ -block-fading,  $L$ -length,  $S$ -sparse channel** described earlier, we established [1] that

1. In the high-SNR regime, the **ergodic capacity** obeys

$$C_{\text{sparse}}(\text{SNR}) = \frac{N - S}{N} \log(\text{SNR}) + O(1).$$

2. To **achieve** the prelog factor  $R_{\text{sparse}} = \frac{N-S}{N}$ , it suffices to use
  - pilot-aided OFDM (with  $N$  subcarriers, of which  $S$  are pilots)
  - with *joint* channel estimation and data decoding.

Key points:

- The effect of *unknown channel support* manifests only in the  $O(1)$  offset.
- Standard non-sparse-channel methods would use  $L$  pilots.
- “Compressed channel sensing” would use  $S$  polylog  $N$  pilots.

[1] A. Pachai-Kannu and P. Schniter, “On communication over unknown sparse frequency selective block-fading channels,” *IEEE Trans. Info. Thy.*, Oct. 2011.

## Practical Communication over the unknown Sparse Channel:

We now propose a communication scheme that...

- is practical, with decode complexity  $O(N \log_2 N + N|S|)$  per  $N$ -block,
- delivers outage rates **matching the optimal prelog factor**  $R_{\text{sparse}} = \frac{N-S}{N}$ ,
- significantly outperforms “compressed channel sensing” (CCS) schemes.

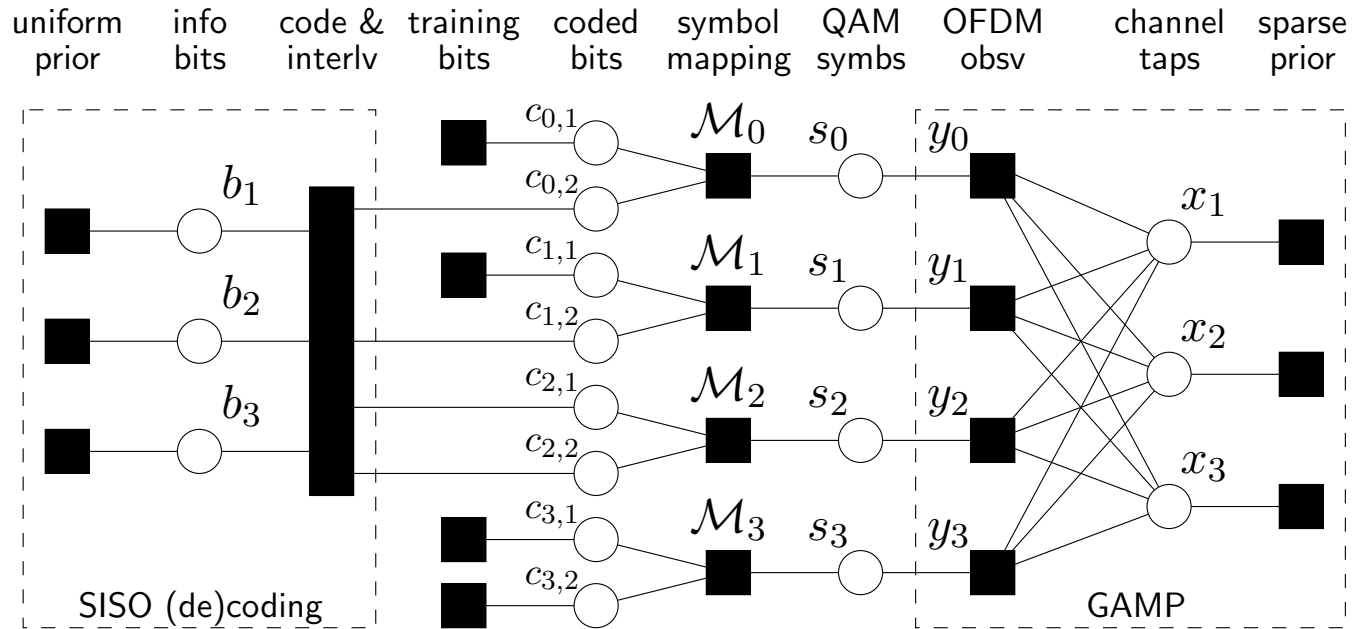
Our scheme uses...

- a conventional transmitter: pilot-aided BICM OFDM,
- a novel receiver: based on belief propagation with the **generalized approximate message passing (GAMP)** algorithm [3] used in a “turbo” configuration [2].

[2] P. Schniter, “Turbo reconstruction of structured sparse signals,” *CISS* 2010.

[3] S. Rangan, “Generalized approximate message passing for estimation with random linear mixing,” *arXiv:1010.5141*, 2010.

### Factor Graph for pilot-aided BICM-OFDM:



○ = random variable

■ = posterior factor

*To jointly infer all random variables, we perform loopy-BP via the sum-product algorithm, using AMP approximations in the GAMP sub-graph.*

## Numerical Results — Perfectly Sparse Channel:

Transmitter:

- LDPC codewords with length  $\sim 10000$  bits.
- $2^M$ -QAM with  $2^M \in \{4, 16, 64, 256\}$  and multi-level Gray mapping.
- OFDM with  $N = 1024$  subcarriers.
- $P$  pilot subcarriers and/or  $T$  training MSBs.

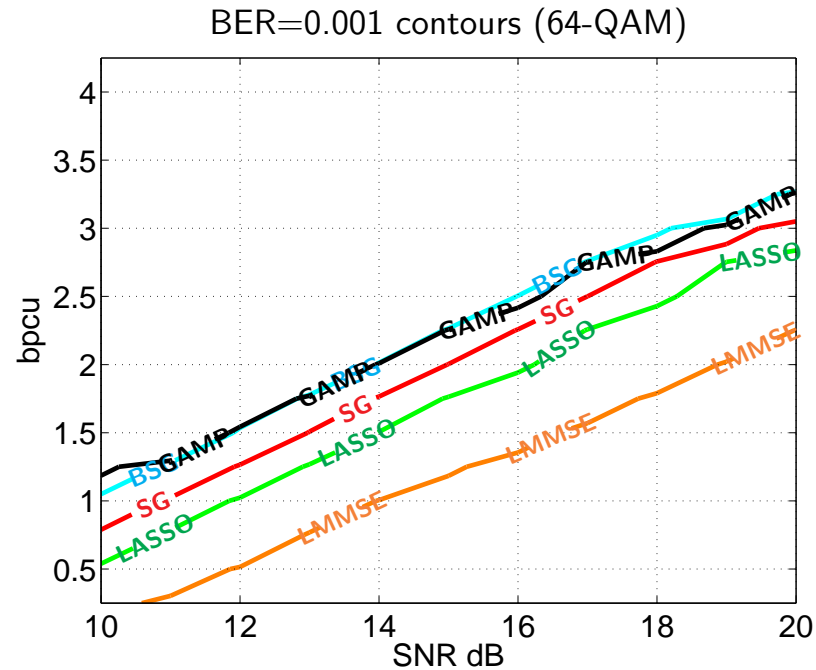
Channel:

- Length  $L = 256 = N/4$ .
- Sparsity  $S = 64 = N/16$ .

Reference Schemes:

- **Pilot-aided LASSO** (i.e., compressed channel sensing) with oracle tuning.
- **Pilot-aided LMMSE**, **support-aware MMSE**, and **info-bit+support-aware MMSE** channel estimates were also tested.

## BER & Outage vs SNR (with $P=L$ pilots & $T=0$ training MSBs):

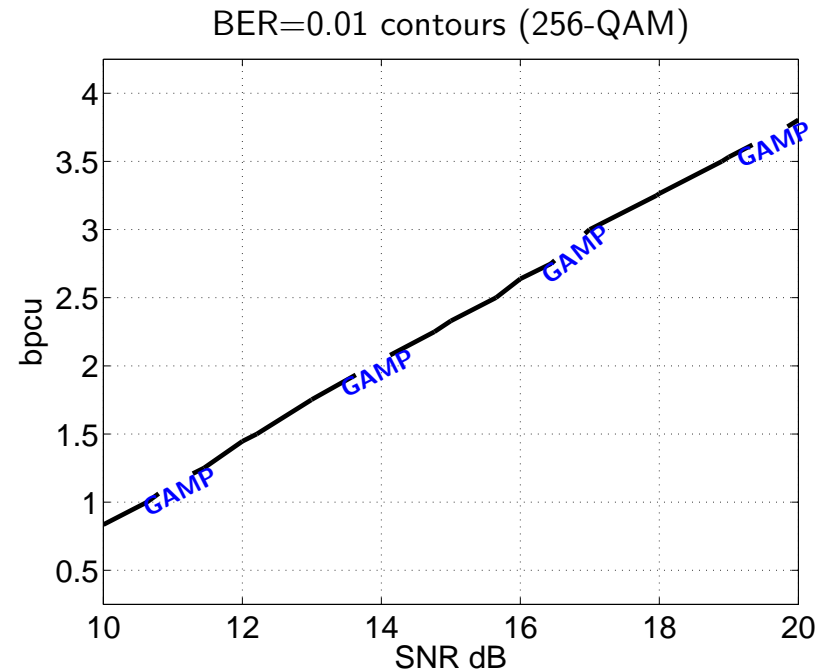
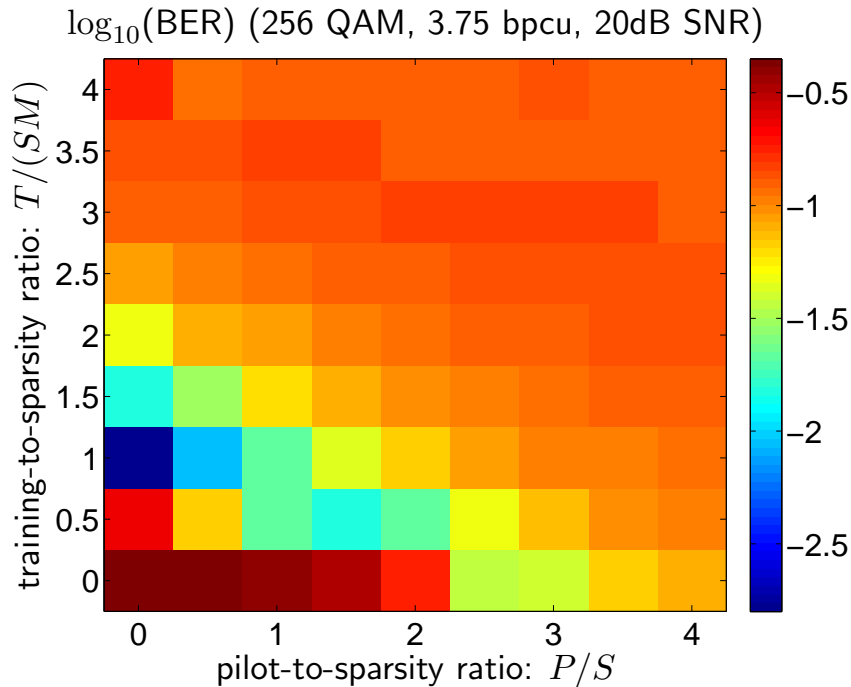


Key points:

- GAMP outperforms both LASSO and the support genie (SG).
- GAMP performs nearly as well as the info-bit+support-aware genie (BSG).
- With  $P = L$ , all approaches yield prelog factor  $R = \frac{N-L}{N} = \frac{3}{4}$ , which falls short of the optimal  $R_{\text{sparse}} = \frac{N-S}{N} = \frac{15}{16}$ .



## BER & Outage vs SNR (with $P=0$ pilots & $T=SM$ training MSBs):

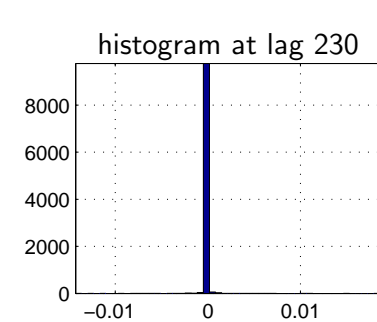
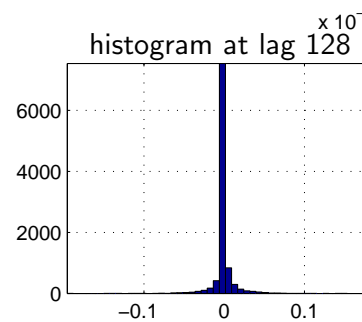
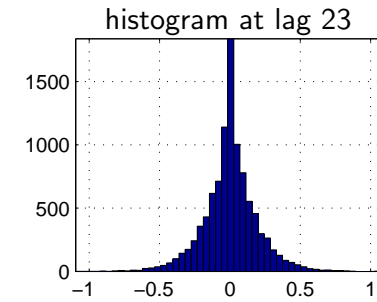
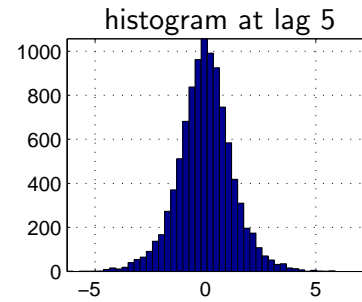
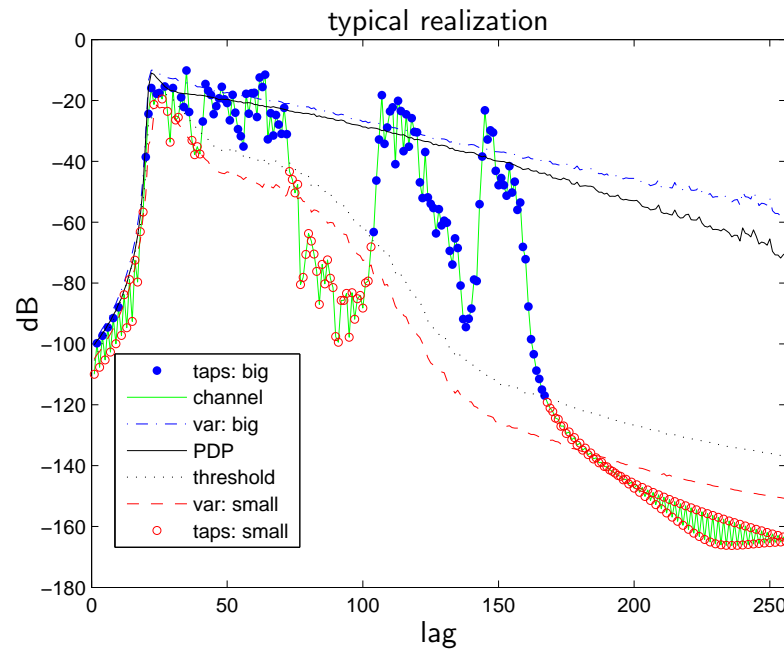


### Key points:

- GAMP favors  $P=0$  pilot subcarriers and  $T=SM$  training MSBs.
  - Precisely the necc/suff redundancy of the capacity-maximizing system!
- GAMP achieves the sparse-channel's capacity-prelog factor,  $R_{\text{sparse}} = \frac{N-S}{N}$ .

## In practice, channel taps are not perfectly sparse, nor i.i.d:

- For example, consider channel taps  $\mathbf{x} = [x_0, \dots, x_{L-1}]$ , where
  - $x_n = x(nT)$  for bandwidth  $T^{-1} = 256$  MHz,
  - $x(t) = h(t) * p_{RC}(t)$ , and
  - $h(t)$  is generated randomly using 802.15.4a outdoor NLOS specs.



- The tap distribution *varies as the lag increases*, becoming more heavy-tailed.
- The **big** taps are *clustered together* in lag, as are the **small** ones.

## Proposed channel model:

- Saleh-Valenzuela (e.g., 802.15.4a) models are accurate but difficult to exploit in receiver design.
- We propose a structured-sparse channel model based on a **2-state Gaussian Mixture** model with **discrete-Markov-chain** structure on the state:

$$p(x_j | d_j) = \begin{cases} \mathcal{CN}(x_j; 0, \mu_j^0) & \text{if } d_j = 0 \text{ "small"} \\ \mathcal{CN}(x_j; 0, \mu_j^1) & \text{if } d_j = 1 \text{ "big"} \end{cases}$$

$$\Pr\{d_{j+1} = 1\} = p_j^{10} \Pr\{d_j = 0\} + (1 - p_j^{01}) \Pr\{d_j = 1\}$$

- Our model is parameterized by the lag-dependent quantities:

$\{\mu_j^1\}$  : big-state power-delay profile

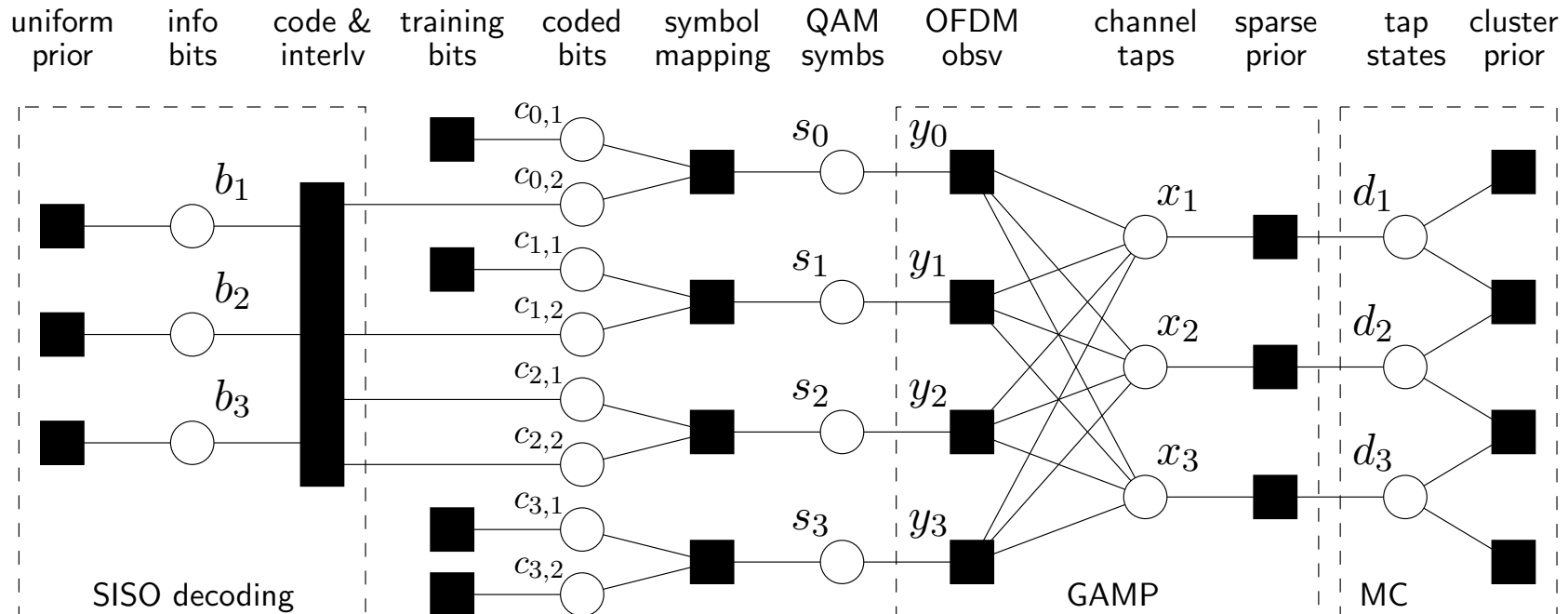
$\{\mu_j^0\}$  : small-state power-delay profile

$\{p_j^{01}\}$  : big-to-small transition probabilities

$\{p_j^{10}\}$  : small-to-big transition probabilities

- Can learn these statistical params from observed realizations via the EM alg.

### Factor graph for pilot-aided BICM-OFDM:



○ = random variable

■ = posterior factor

*To jointly infer all random variables, we perform loopy-BP via the sum-product algorithm, using AMP approximations in the GAMP sub-graph.*

## Numerical results:

### Transmitter:

- OFDM with  $N = 1024$  subcarriers.
- 16-QAM with multi-level Gray mapping
- LDPC codewords with length  $\sim 10000$  yielding spectral efficiency of 2 bpcu.
- $P$  pilot subcarriers and  $T$  training MSBs.

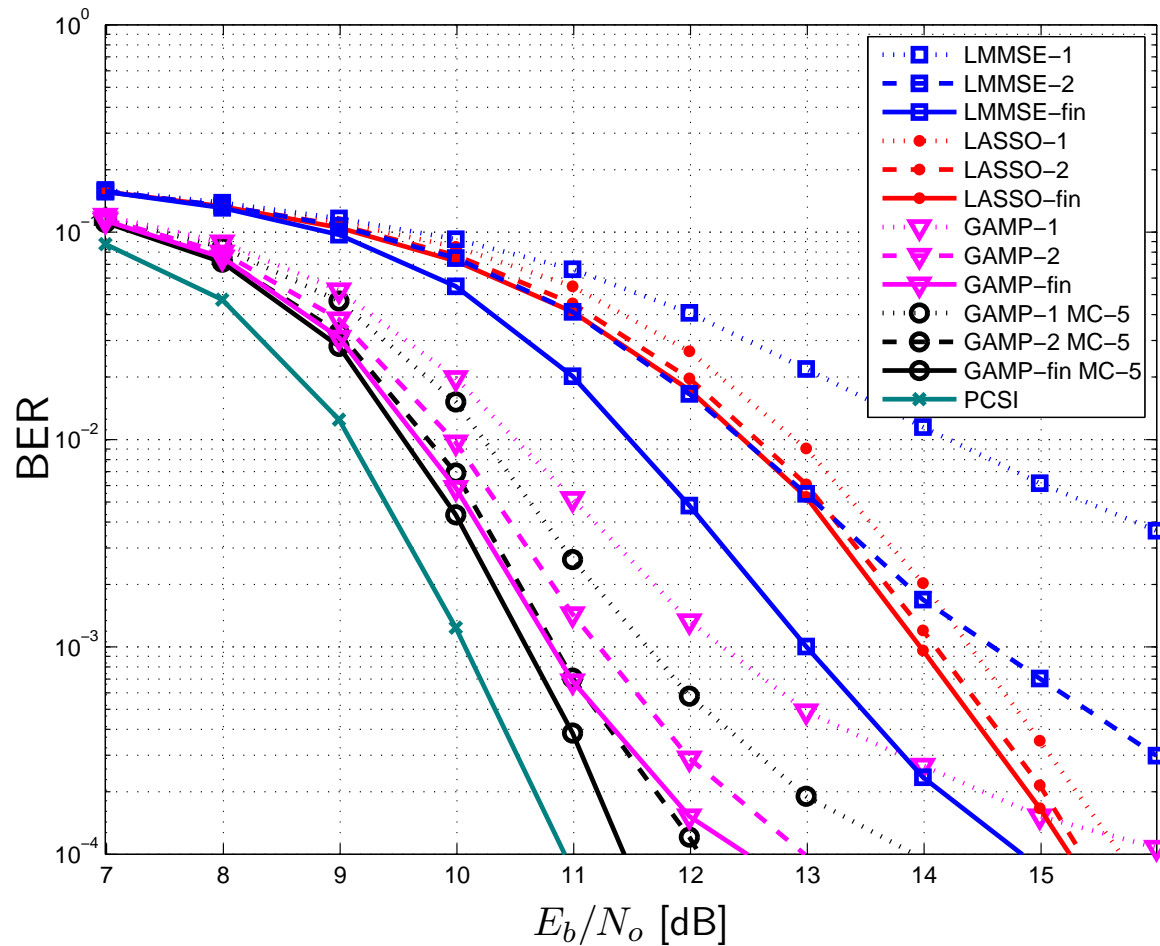
### Channel:

- 802.15.4a outdoor-NLOS (not our Gaussian-mixture model!)
- Length  $L = 256 = N/4$ .

### Reference Channel Estimation / Equalization Schemes:

- soft-input soft-output (SISO) versions of LMMSE and LASSO.
- perfect-CSI genie.

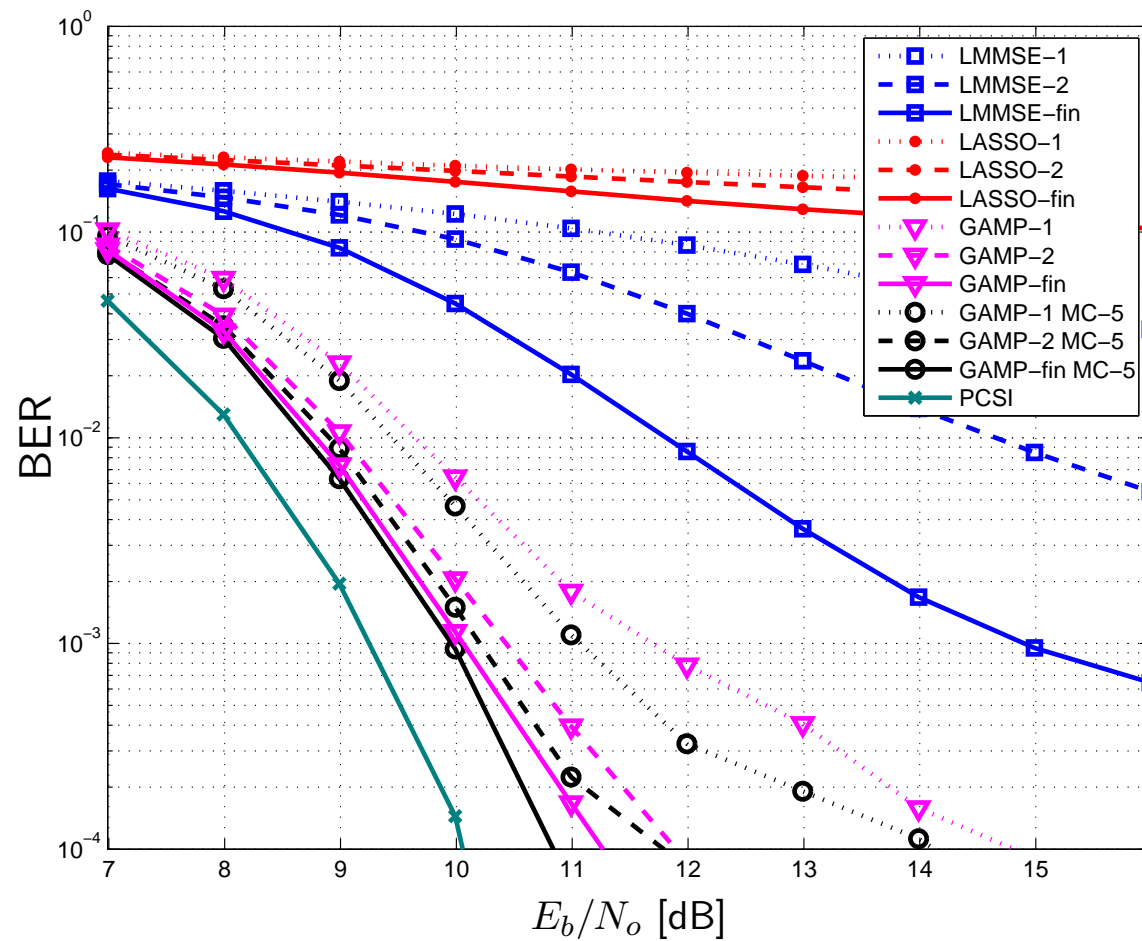
## BER versus $E_b/N_o$ for $P = 224$ pilots and $T = 0$ training MSBs:



Note 4dB improvement over (turbo) LASSO.

Only 0.5dB from perfect-CSI genie!

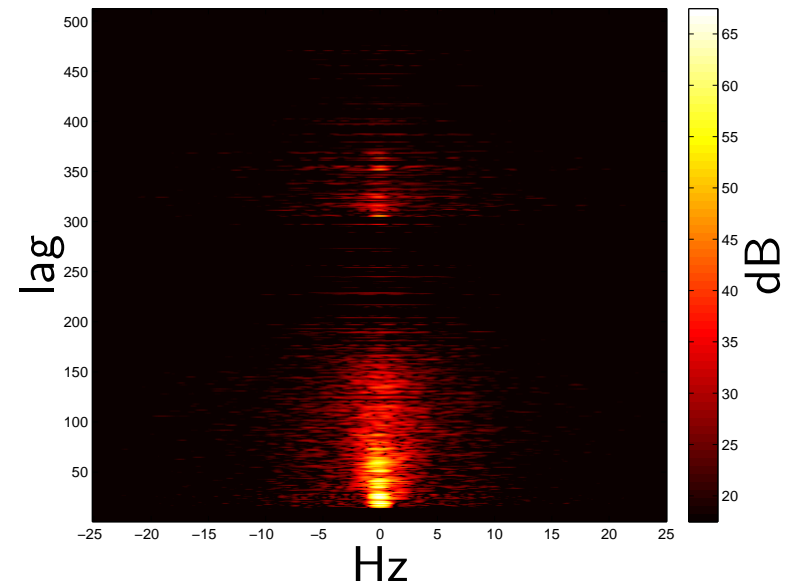
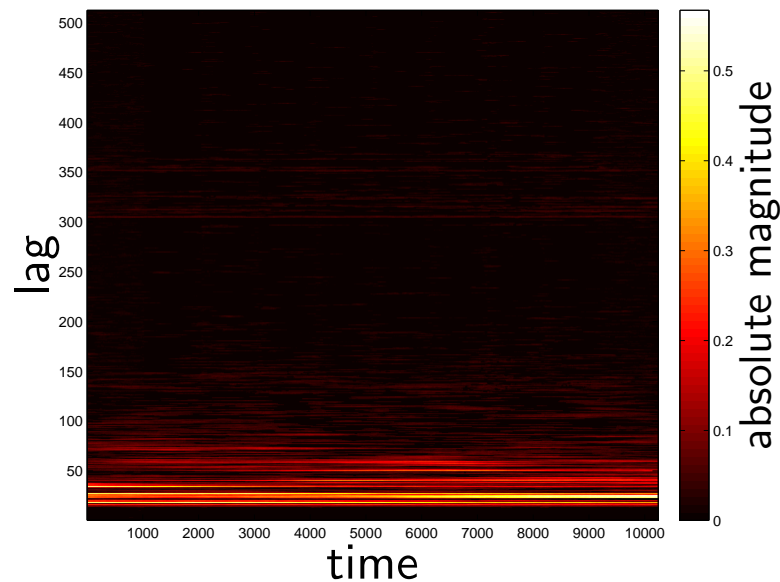
## BER versus $E_b/N_o$ for $P = 0$ pilots and $T = 448$ training MSBs:



Use of training MSBs gives 1dB improvement over use of pilot subcarriers!

## Communications over Underwater Channels:

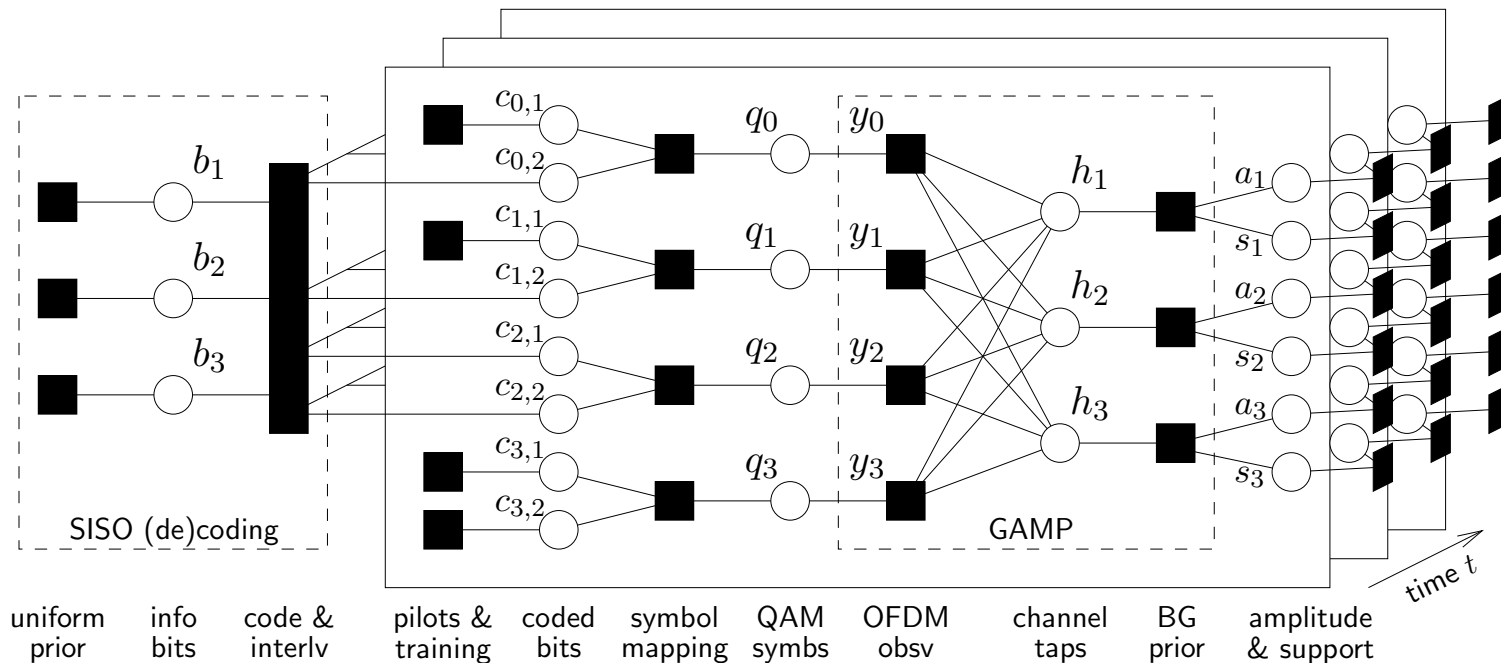
- SPACE-08 Underwater Experiment 2920156F038\_C0\_S6
- Time-varying channel response estimated using WHOI M-sequence:



- The channel is nearly over-spread:  $f_d T_s L = 20 \times \frac{1}{10000} \times 400 = 0.8$  !
- Can't afford to ignore structure of temporal variations!



## BICM-OFDM Factor Graph with Temporal Channel Structure:



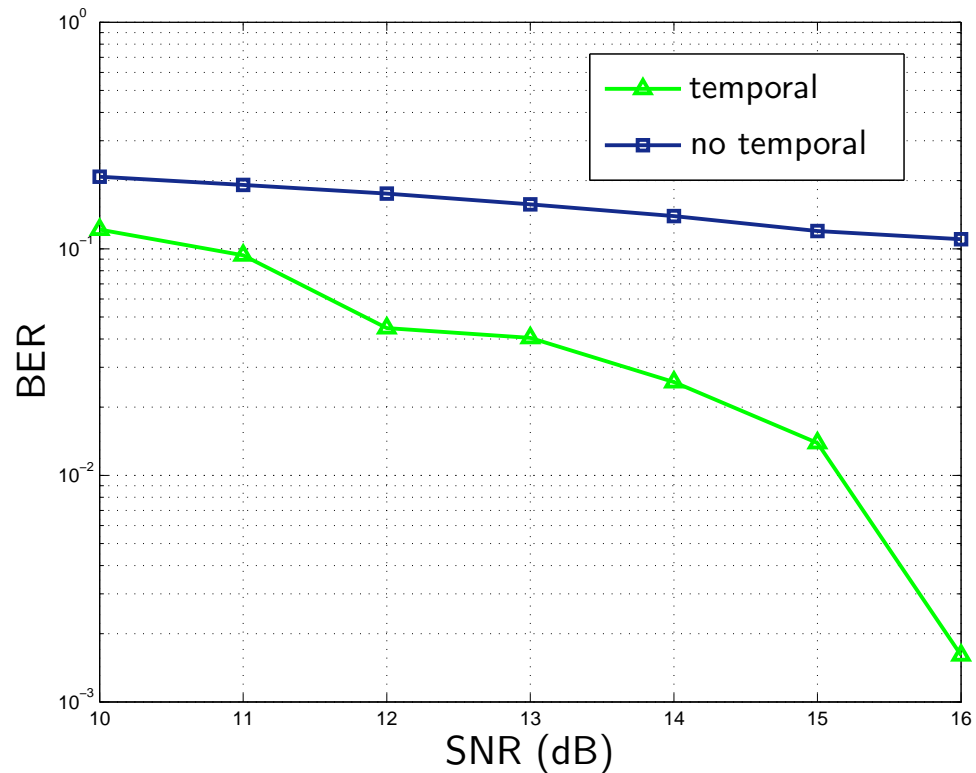
- Channel taps are modeled as independent Bernoulli-Gaussian processes:
  - each tap's amplitude follows a [temporal Gauss-Markov chain](#)
  - each tap's on/off state follows a [temporal discrete-Markov chain](#)

[4] P. Schniter and D. Meng, "A Message-Passing Receiver for BICM-OFDM over Unknown Time-Varying Sparse Channels," *Allerton* 2011.

## Performance versus SNR:

### Settings:

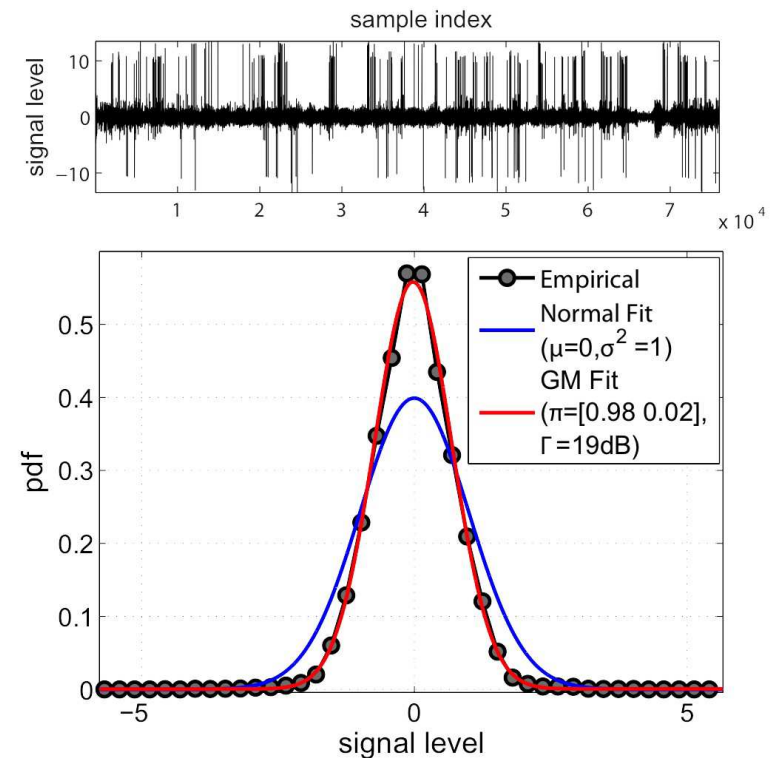
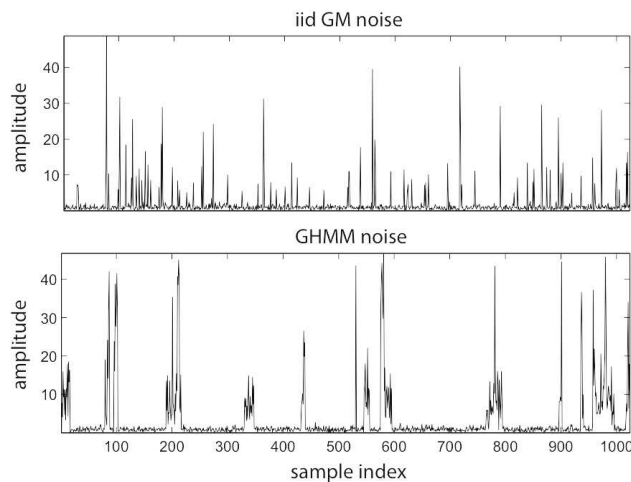
- experimentally measured underwater channel
- 16-QAM
- 1024 total tones
- 0 pilot tones
- 256 training MSBs
- LDPC length 10k
- LDPC rate 0.5



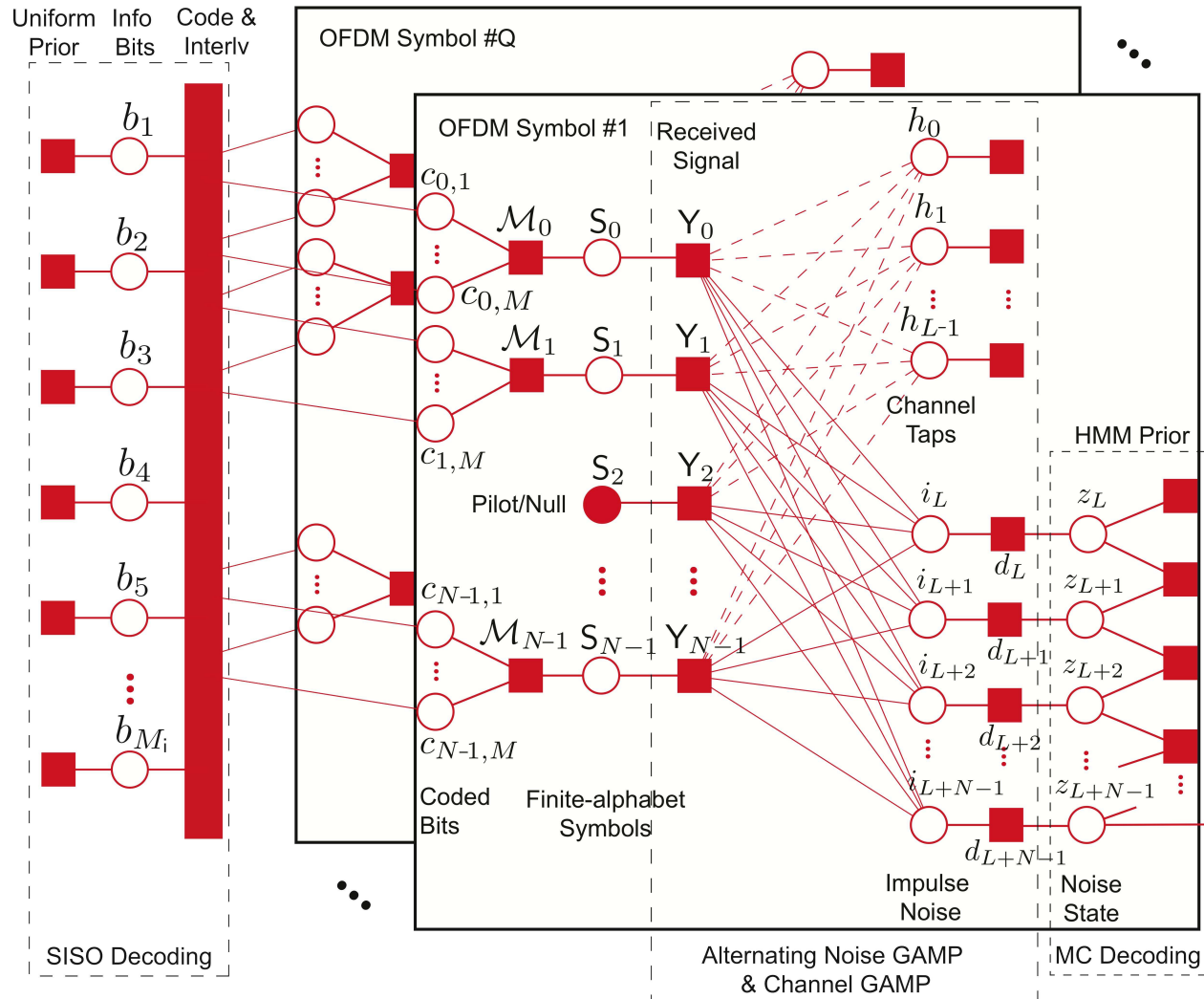
Exploiting the [persistence in channel support and channel amplitudes](#) was critical in this difficult underwater application.

## Communications in Impulsive Noise:

- In many wireless and power-line communication systems, the (time-domain) noise is not Gaussian but **impulsive**.
- The marginal noise statistics are well captured by a **2-state Gaussian mixture** (i.e., Middleton class-A) model.
- Noise burstiness is well captured by a **discrete Markov chain** on the noise state.



# Factor Graph for pilot-aided BICM-OFDM:

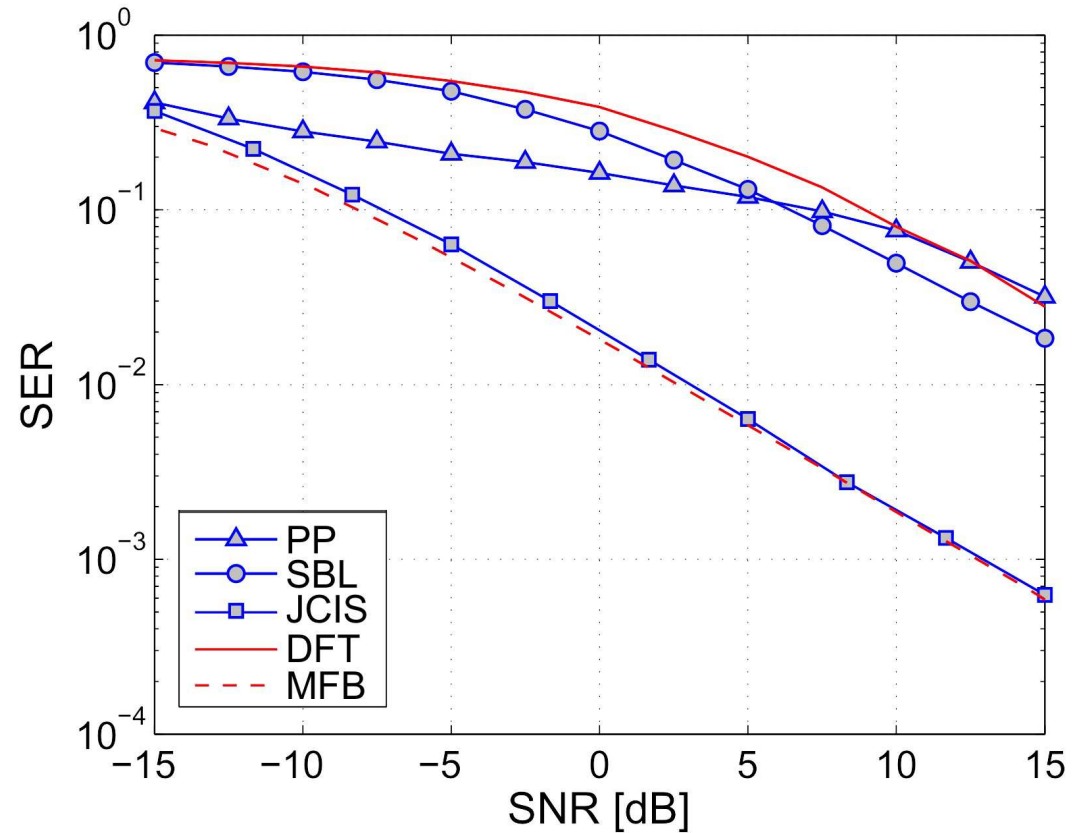


[5] M. Nassar, P. Schniter, and B. Evans, "A Factor-Graph Approach to Joint OFDM Channel Estimation and Decoding in Impulsive Noise Environments," *IEEE Trans. Signal Process.*, 2014.

## Numerical Results — Uncoded Case:

Settings:

- 5 channel taps
- GM noise
- 256 total tones
- 15 pilot tones
- 80 null tones
- 4-QAM

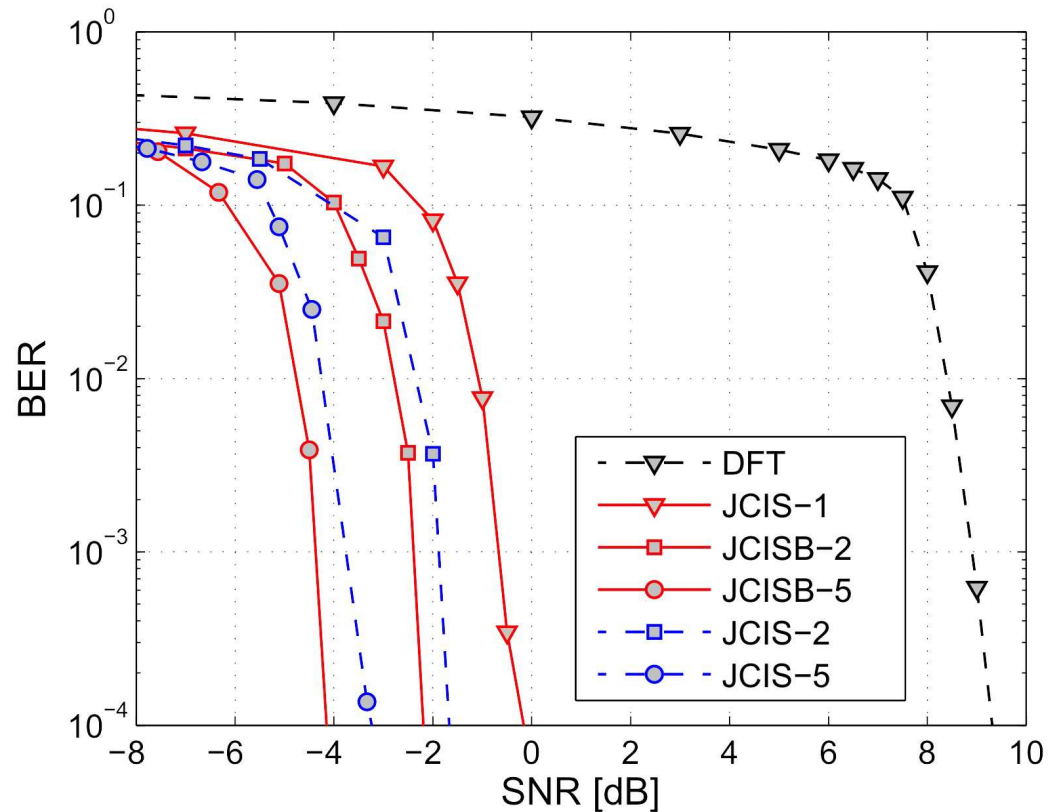


Proposed “joint channel/impulsive-noise/symbol” estimation (JCIS) scheme gives  $\sim 15$  dB gain over previous state-of-the-art and performs within 1 dB of MFB!

## Numerical Results — Coded Case:

### Settings:

- 10 channel taps
- GM noise
- 1024 total tones
- 150 pilot tones
- 0 null tones
- 16-QAM
- LDPC
- Rate 0.5
- Length 60k



Proposed “joint channel/impulsive-noise/symbol/bit” estimation (JCISB) scheme gives  $\sim 15$  dB gain over traditional DFT-based receiver!

## Conclusions:

- At wide bandwidths, channel **impulse responses** are approximately **sparse**.
  - Sparsity increases the pre-log factor of high-SNR noncoherent ergodic capacity.
  - AMP-based joint channel-estimation/decoding delivers outage rates that empirically match the capacity pre-log factor.
  - Channels impulses are in fact structured-sparse, and exploiting this structure leads to additional performance gains.
  - Sparsity can also be exploited in time-varying channels.
- **Impulsive noise** is another source of sparsity in communications.
  - AMP-based joint channel-estimation/impulse-estimation/decoding delivers error-rates that approach the matched-filter bound.